

# Seminar of Computer Networks

## Homework 1

You can do the homeworks alone or in groups of two students. Write a single solution with both names.

Make sure that the solutions are typewritten or clear to read.

**Hand in your solutions and keep a copy for yourself.** After the due date we will post the solutions and in the final exam we will ask you to explain us what were your mistakes.

**Due date:** 18/5/2011, before the class.

**Problem 1.** Let  $X$  be a discrete random variable that takes values different than 0 with expectation  $\mu = \mathbf{E}[X]$ . Define the random variable  $Y = 1/X$ . Does it hold in general that  $\mathbf{E}[Y] = 1/\mu$ ? You should provide a simple proof if you believe that it is true, or a simple counterexample if you believe that it isn't.

**Problem 2.** Consider a random graph  $G = (V, E)$  with  $|V| = n$  nodes created according to the Erdős-Rényi  $G_{n,p}$  random-graph model. Let us define a *3-star* to be a subgraph  $V' = \{v_0, v_1, v_2, v_3\} \subseteq V$  of  $V$  such that all the edges  $(v_0, v_i)$  for  $i = 1, 2, 3$  exist in  $G$ , and none other of the edges  $(v_1, v_2)$ ,  $(v_1, v_3)$ , and  $(v_2, v_3)$  exist in the graph (but other edges may exist, including edges between nodes in  $V'$  and other nodes in  $V$ ).

1. What is the expected number of 3-stars in  $G$ ?
2. Define similarly a *k-star* (a subset  $V' = \{v_0, v_1, \dots, v_k\} \subseteq V$  such that for  $i = 1, \dots, k$  we have that  $(v_0, v_i) \in E$ , and for  $i, j = 1, \dots, k$  we have that  $(v_i, v_j) \notin E$ ). What is the expected number of *k-stars* in  $G$ ?

**Problem 3.** Consider the following modification of the Barabassi-Albert preferential attachment model that we did in class: When a new node arrives at time  $t$  again it comes with  $\ell$  edges. However, this time each edge selects a node  $v$  with probability proportional to the degree  $d_v$  plus a constant  $c$ , that is, the probability equals

$$\frac{d_v + c}{(t-1)(2\ell + c)},$$

where  $c \geq -\ell$ , as we describe at the end of Chapter 4 in the notes (so for  $c = 0$  this is the Barabassi-Albert model). Show that the degree distribution that we obtain as  $t \rightarrow \infty$  is approximately a power law with exponent  $3 + c/\ell$ .

**Problem 4.** Let  $G = (V, E)$  be the circle graph, that is, the graph with node set  $V = \{v_1, v_2, \dots, v_n\}$  and edge set the  $E = \{(v_i, v_{i+1}); i = 1, \dots, n\}$  (where we assume that the edge  $(v_n, v_{n+1})$  stands for edge  $(v_n, v_1)$ ). Find the best partitioning for  $G$  that maximizes the modularity (and show that it is the best possible).

**Note:** Use the following fact without having to prove it: Among all the possible sets of  $n$  numbers  $a_1, a_2, \dots, a_n$  for which we know that the sum  $\sum_{i=1}^n a_i = L$  is some fixed value  $L$ , the sum of the squares  $\sum_{i=1}^n a_i^2$  is minimized when we have that  $a_1 = a_2 = \dots = a_n = L/n$ .

**Problem 5.** Write a simple program in the programming language that you prefer to calculate the degree distributions of the graphs in

<http://www.dis.uniroma1.it/~socialnet/graphs>

and plot them in regular and in log-log scale. For plotting you can also use any program that you wish (Excel, OpenOffice, Matlab, R, gnuplot, etc.).

Next find the best power-law function that fits the data. To do that there are various definitions of “best.” A simple one (not necessarily the best) is the following. Recall that a power law distribution function is a distribution function of the form

$$y(x) = b \cdot x^a.$$

If we take logarithms in both sides we have

$$\ln y(x) = \ln b + a \ln x,$$

that is, we have a linear relationship between  $\ln x$  and  $\ln y$ . We can find the parameters  $a$  and  $b$  that create the line that best fits a set of data  $(x_i, y_i)$  by minimizing the mean-square error of the logarithms. That is, we find the values of  $a$  and  $b$  that minimize the expression

$$\text{RSS} = \sum_{i=1}^N (\ln y_i - (a \ln x_i + \ln b))^2,$$

where  $(x_i, y_i)$  is an input pair, and  $N$  the total number of points. In our case a pair  $(x_i, y_i)$  corresponds to a pair (degree, number of nodes with that degree), and  $N$  is the number of different degrees. This would give the line that best fits the points in the log-log plot as measured by the *residual sum of squares* (RSS) error.

Since the relationship between the  $\ln y$  and  $\ln x$  is linear, it turns out that the values of  $a$  and  $\ln b$  that minimize this expression are given by

$$a = \frac{S_{xy}}{S_{xx}} \quad \text{and} \quad \ln b = \widehat{\ln y} - a \widehat{\ln x},$$

where

$$\begin{aligned} \widehat{\ln x} &= \frac{1}{N} \sum_{i=1}^N \ln x_i \\ \widehat{\ln y} &= \frac{1}{N} \sum_{i=1}^N \ln y_i \\ S_{xx} &= \sum_{i=1}^N (\ln x_i - \widehat{\ln x})^2 \\ S_{xy} &= \sum_{i=1}^N (\ln x_i - \widehat{\ln x})(\ln y_i - \widehat{\ln y}). \end{aligned}$$

After you compute the parameters  $a$  and  $b$ , plot in the same figures with the points  $(x_i, y_i)$  that you plotted earlier the function  $y = b \cdot x^a$ .

Mail the code that you write to compute the degree distributions [aris@cs.brown.edu](mailto:aris@cs.brown.edu), with subject: “Seminar Computer Networks - Homework 1.”